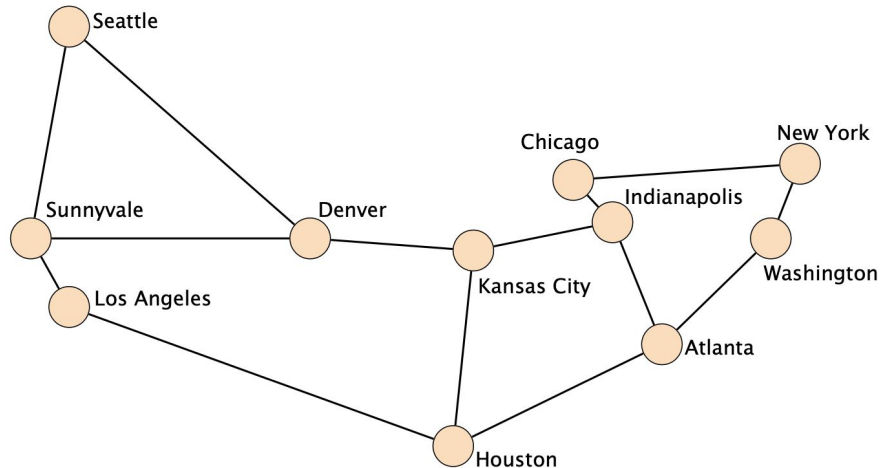


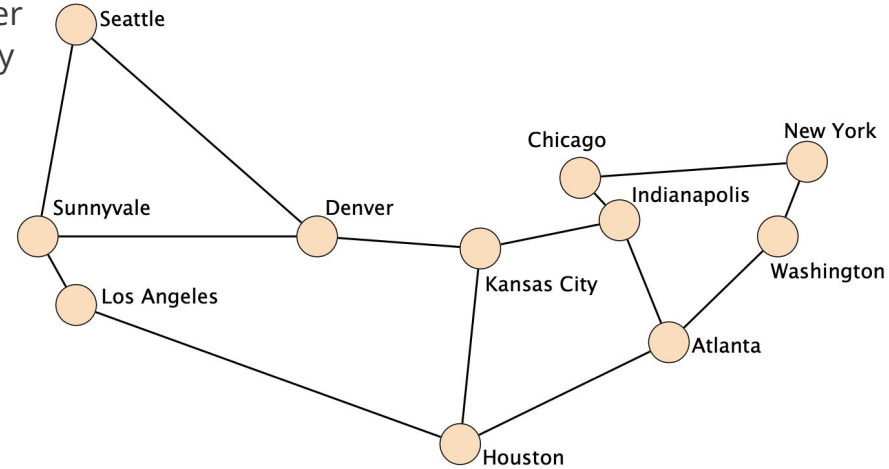
# Link Weight Configurations

- The Abilene network was a high-performance backbone network in the US. You are the network operator in charge and you have to configure the link weights in the network. Initially, all links have a weight of one and routers will always use the shortest-path available to reach a destination.



# Link Weight Configurations

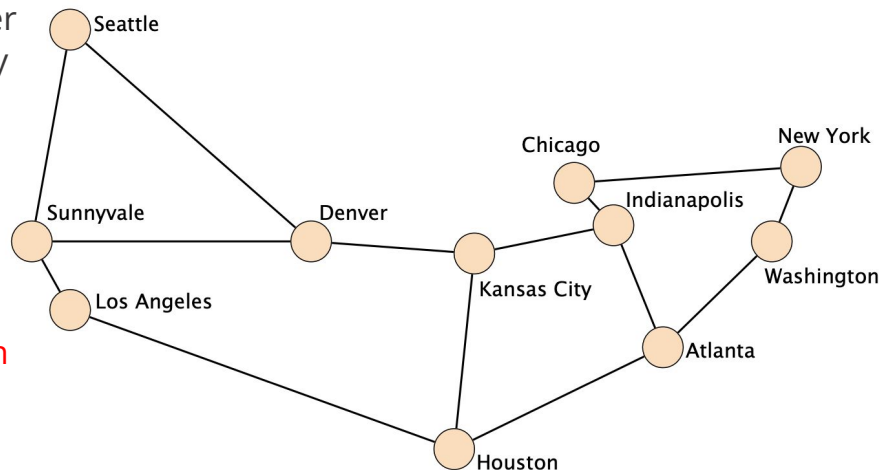
Is it possible to configure the link weights such that the packets sent by the router located in Los Angeles to the router located in New York follow one path while the packets sent by the router located in New York to the router located in Los Angeles follow a completely different path?



# Link Weight Configurations

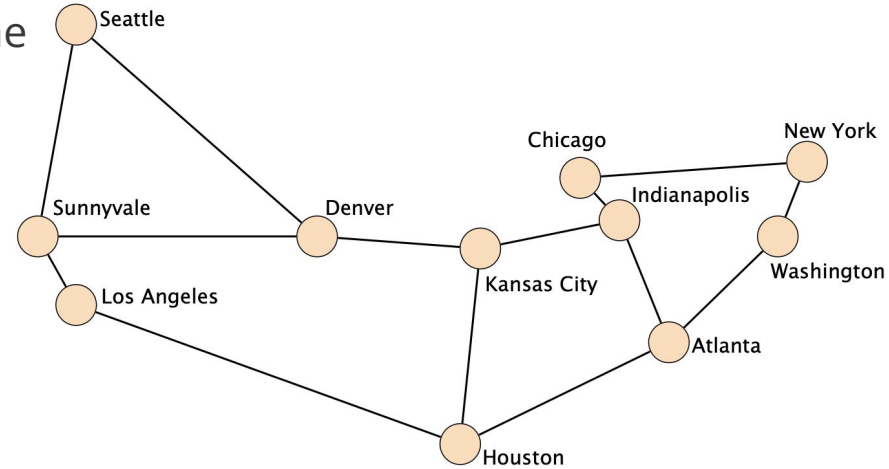
Is it possible to configure the link weights such that the packets sent by the router located in Los Angeles to the router located in New York follow one path while the packets sent by the router located in New York to the router located in Los Angeles follow a completely different path?

**Solution: Not possible.** We consider only links which have the same weight in both directions. If the two routers would use different paths for the two traffic directions, the two paths would need different total weights. That implies that one path is shorter and one router is not using the shortest-path available. A contradiction to our initial assumption.



# Link Weight Configurations

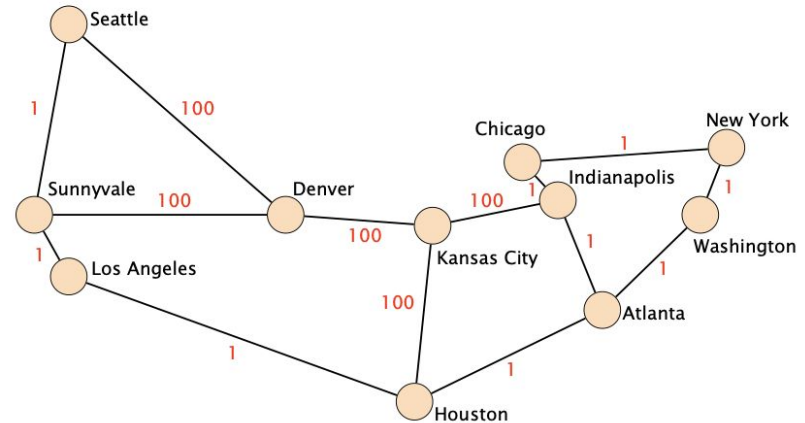
Assume that the routers located in Denver and Kansas City need to exchange lots of data on the direct link. Can you configure the link weights such that the link between these two routers does not carry any packet sent by any other router in the network?



# Link Weight Configurations

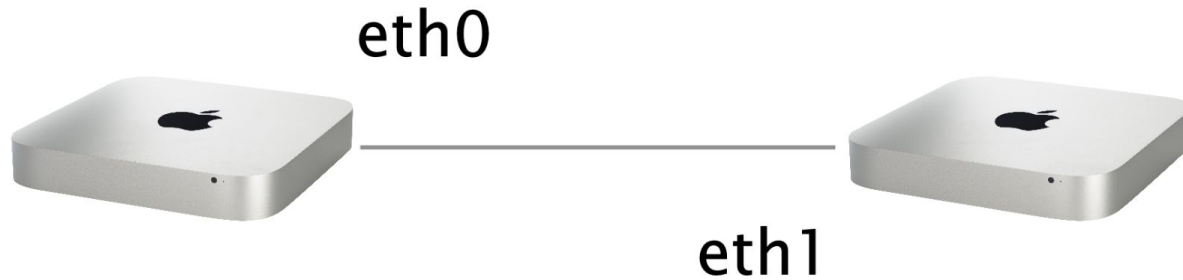
Assume that the routers located in Denver and Kansas City need to exchange lots of data on the direct link. Can you configure the link weights such that the link between these two routers does not carry any packet sent by any other router in the network?

**Solution:**



# Link Layer

How do **local** computers communicate?



# Link Layer

1. What is a link?
2. How do we share a network medium?
3. How do we identify link adapters?
4. What is Ethernet?
5. How do we interconnect segments at the link layer?



# Link Layer

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# Link Layer

Link

Communication  
medium

and

Network  
adapter



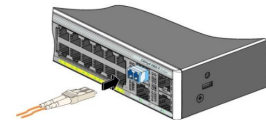
Wifi



Ethernet



Fiber



# Link Layer

terminology:

- hosts and routers: **nodes**
- communication channels that connect adjacent nodes along communication path: **links**
  - wired
  - wireless
  - LANs
- layer-2 packet: *frame*, encapsulates datagram

**Link layer** has responsibility of transferring datagram from one node to **physically adjacent** node over a link

# Link Layer

- datagram transferred by different link protocols over different links:
  - e.g., WiFi on first link, Ethernet on next link
- each link protocol provides different services
  - e.g., may or may not provide reliable data transfer over link

transportation analogy:

- trip from Princeton to Lausanne
  - limo: Princeton to JFK
  - plane: JFK to Geneva
  - train: Geneva to Lausanne
- tourist = datagram
- transport segment = communication link
- transportation mode = link-layer protocol
- travel agent = routing algorithm

# Link Layer Services

- framing, link access:
  - encapsulate datagram into frame, adding header, trailer
  - channel access if shared medium
  - “MAC” addresses in frame headers identify source, destination
- reliable delivery between adjacent nodes
  - seldom used on low bit-error links
  - wireless links: high error rates
  - **Question:** why both link-level and end-end reliability?

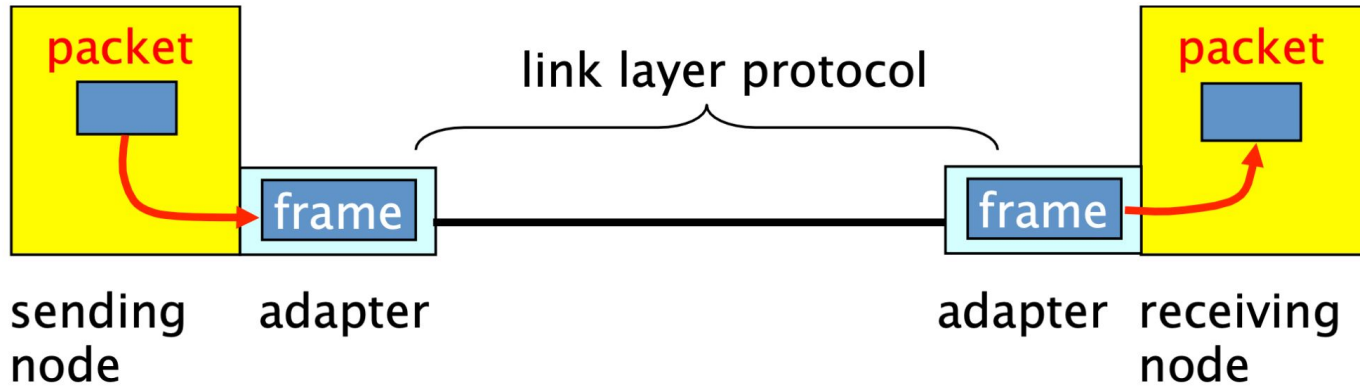
# Link Layer Services

- flow control:
  - pacing between adjacent sending and receiving nodes
- error detection:
  - errors caused by signal attenuation, noise.
  - receiver detects errors, signals retransmission, or drops frame
- error correction:
  - receiver identifies and corrects bit error(s) without retransmission
- half-duplex and full-duplex:
  - with half duplex, nodes at both ends of link can transmit, but not at same time

# Where is the Link Layer Implemented?

- each-and-every host
- link layer implemented in *network interface card* (NIC) or on a chip
  - Ethernet, WiFi card or chip
  - implements link, physical layer
- attaches into host's system buses
- combination of hardware, software, firmware

# Network Adapters Communicate via the Medium



sending side:

- encapsulates datagram in frame
- adds error checking bits, reliable data transfer, flow control, etc.

receiving side:

- looks for errors, reliable data transfer, flow control, etc.
- extracts datagram, passes to upper layer at receiving side



# Link Layer

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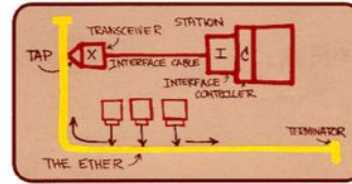
# Some Mediums are Multi-Access (>1 Host can Communicate at a Time)



Wireless  
networks



Satellite  
networks



original Ethernet  
networks



Cellular  
networks

# Multi-Access Problem - Collisions

## Problem

collisions lead  
to garbled data

## Solution

distributed algorithm  
for sharing the channel

When can each node transmit?

# There are Effectively Three Ways to Manage Multiple Access

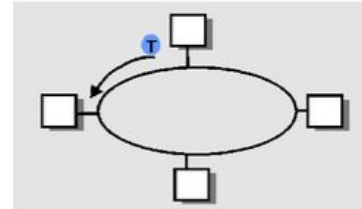
**Divide the channel into pieces**

either in time or in frequency



**Take turns**

pass a token for the right to transmit



**Random access**

allow collisions, detect them and then recover

# There are Effectively Three Ways to Manage Multiple Access

channel partitioning, by time, frequency or code

- Time Division, Frequency Division

taking turns

- polling from central site, token passing
- Bluetooth, FDDI, token ring

random access (dynamic),

- ALOHA, S-ALOHA, CSMA, CSMA/CD
- carrier sensing: easy in some technologies (wire), hard in others (wireless)
- CSMA/CD used in Ethernet
- CSMA/CA used in 802.11

# Carrier Sense Multiple Access

simple CSMA: listen before transmit:

- if channel sensed idle: transmit entire frame
- if channel sensed busy: defer transmission

human analogy: don't interrupt others

CSMA/CD: CSMA with collision detection

- collisions detected within short time
- colliding transmissions aborted, reducing channel wastage
- collision detection easy in wired, difficult with wireless

human analogy: the polite conversationalist

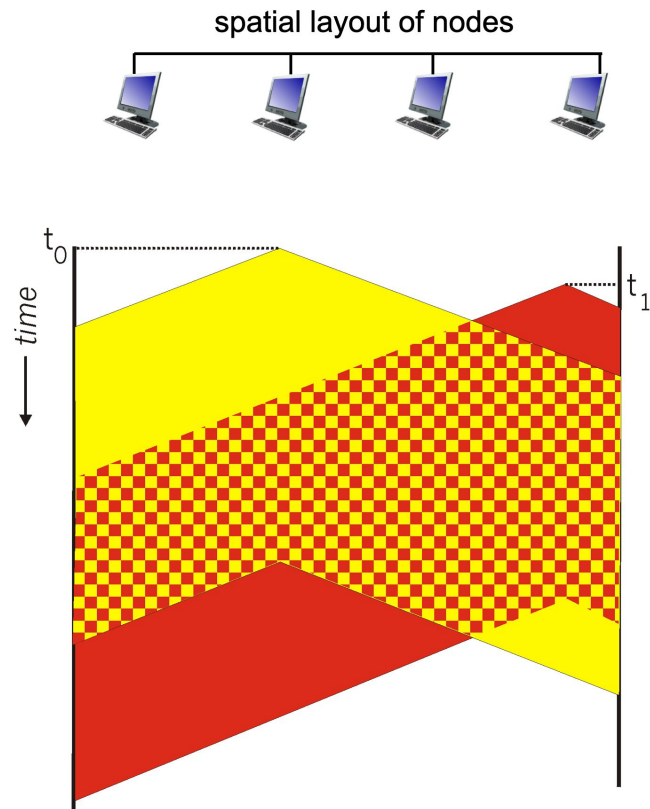
# Pure CSMA

collisions can still occur with carrier sensing:

- propagation delay means two nodes may not hear each other's just-started transmission

collision: entire packet transmission time wasted

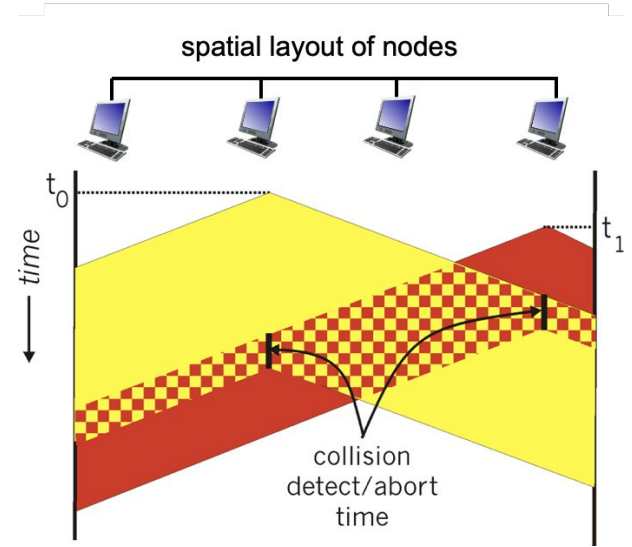
- distance & propagation delay play role in determining collision probability



# CSMA/CD

CSMA/CD reduces the amount of time wasted in collisions

- transmission aborted on collision detection





# Link Layer

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Medium Access Control addresses

# MAC Addresses

identify the sender & receiver adapters

used on a link

are uniquely assigned

hard-coded into the adapter when built

use a flat space of 48 bits

allocated hierarchically

## MAC Addresses are Hierarchical

34 : 36 : 3b : d2 : 8a : 86

# The First 24 bits Represent the Vendor

34:36:3b:d2:8a:86

Apple, Inc.  
1 Infinite Loop  
Cupertino CA 95014  
US

see <http://standards-oui.ieee.org/oui/oui.txt>

# The Second 24 bits are Assigned by the Vendor to the Adapter

34:36:3b:d2:8a:86

assigned by Apple  
to my adapter

If all bits are Set, it's a **Broadcast** Address

ff:ff:ff:ff:ff:ff

enables to send a frame to  
*all* adapters on the link

**By Default, Adapters only Decapsulates Frames Addressed to the Local MAC or the Broadcast Address**



# By Default, Adapters only Decapsulates Frames Addressed to the Local MAC or the Broadcast Address

Workaround is **promiscuous** mode: enables to decapsulate everything, independently of the destination MAC

# Why Do We Need MAC Addresses (and not just use IPs)?

Adapters must be identified during bootstrap  
need to talk to an adapter to give it an IP address

# Network Adapter Bootstrap

Two problems to solve:

1. Who am I?                      How do I acquire an IP address?
  
2. Who are you?                      Given an IP, how do I find which MAC to send to?

# Network Adapter Bootstrap

Two problems to solve:

1. Who am I?                      How do I acquire an IP address?  
**Dynamic Host Configuration Protocol**
  
2. Who are you?                      Given an IP, how do I find which MAC to send to?  
**Address Resolution Protocol**